

Telescopers for differential forms with one parameter*

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Abstract

Telescopers for a function are linear differential (resp. difference) operators annihilating the definite integral (resp. definite sum) of this function. They play a key role in Wilf-Zeilberger theory and algorithms for computing them have been extensively studied in the past thirty years. In this paper, we introduce the notion of telescopers for differential forms with D -finite function coefficients. These telescopers appear in several areas of mathematics, for instance parametrized differential Galois theory and mirror symmetry. We give a sufficient and necessary condition for the existence of telescopers for a differential form and describe a method to compute them if they exist. Algorithms for verifying this condition are also given.

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1 Introduction

In the Wilf-Zeilberger theory, telescopers usually refer to the operators in the output of the method of creative telescoping, which are linear differential (resp. difference) operators annihilating the definite integrals (resp. the definite sums) of the input functions. The telescopers have emerged at least from the work of Euler [17] and have found many applications in the various areas of mathematics such as combinatorics, number theory, knot theory as well as others (see Section 7 of [20] for details). In particular, telescopers for a function are often used to prove the identities involving this function or even obtain a simpler expression for the definite integral or sum of this function. As a clever and algorithmic process for constructing telescopers, creative telescoping firstly appeared as a term in the essay of van der Poorten on Apréy's proof of the irrationality of $\zeta(3)$ [31]. However, it was Zeilberger and his collaborators [3, 29, 36, 37, 40] in the early 1990s who equipped creative telescoping with a concrete meaning and formulated it as an algorithmic tool. Since then, algorithms for creative telescoping have been extensively studied. Based on the techniques used in the algorithms, the existing algorithms are divided into four generations, see [13] for the details. Most recent algorithms are called reduction-based algorithms which were first introduced by Bostan et al. in [6] and further developed in (for example, [7, 8, 14, 15]). The termination of these algorithms relies on the existence of telescopers. The question for which input functions the algorithms will terminate has been answered in [1, 2, 10, 16, 38] for several classes of functions such as rational functions and hypergeometric functions as well as others. The algorithmic framework for creative telescoping is now called the Wilf-Zeilberger theory.

Most of algorithms for creative telescoping focus on the case of one bivariate function as input. There are only a few algorithms which deal with multivariate case (see, for example, [9, 11, 12, 21]). It is still a challenge to develop the multivariate analogue of the existing algorithms (see Section 5 of [13]). In the language of differential forms (with m variables and one parameter), the results in [11] and [21] dealt with the cases of differential 1-forms and differential m -forms respectively. On the other hand, in the applications to other domains such as mirror symmetry (see [23, 26, 27]), one needs to deal with the case of differential p -forms with $1 \leq p \leq m$. Below is an example.

Example 1 *Consider the following one-parameter family of the quintic polynomials*

$$W(t) = \frac{1}{5}(x_1^5 + x_2^5 + x_3^5 + x_4^5 + x_5^5) - tx_1x_2x_3x_4x_5$$

where t is a parameter. Set

$$\omega = \sum_{i=1}^5 \frac{(-1)^{i-1} x_i}{W(t)} dx_1 \wedge \cdots \wedge \widehat{dx}_i \wedge \cdots \wedge dx_5.$$

To obtain the Picard-Fuchs equation for the mirror quintic, the geometers want to compute a fourth order linear differential operator L in t and ∂_t such that $L(\omega) = d\eta$ for some differential 3-form η . Here one has that

$$L = (1 - t^5) \frac{\partial^4}{\partial t^4} - 10t^4 \frac{\partial^3}{\partial t^3} - 25t^3 \frac{\partial^2}{\partial t^2} - 15t^2 \frac{\partial}{\partial t} - 1.$$

Set $\theta_t = t\partial/\partial t$. Then

$$\tilde{L} = -\frac{1}{5^4} L \frac{1}{t} = \theta_t^4 - 5t(5\theta_t + 1)(5\theta_t + 2)(5\theta_t + 3)(5\theta_t + 4)$$

and the equation $\tilde{L}(y) = 0$ is the required Picard-Fuchs equation.

We call the operator L appearing in the above example a telescope for the differential form ω (see Definition 4). In this paper, we study the telescopes for differential forms with D -finite function coefficients. Instead of the geometric method used in [23, 26, 27], we provide an algebraic treatment. We give a sufficient and necessary condition guaranteeing the existence of telescopes and describe a method to compute them if they exist. In addition, we also present algorithms to verify this condition.

The rest of this paper is organized as follows. In Section 2, we recall differential forms with D -finite function coefficients and introduce the notion of telescopes for differential forms. In Section 3, we give a sufficient and necessary condition for the existence of telescopes, which can be considered as a parametrized version of Poincaré lemma on differential manifolds. In Section 4, we give two algorithms for verifying the condition presented in Section 3.

Notations: The following notations will be frequently used throughout this paper.

$$\begin{aligned} \partial_t: & \text{ the usual derivation } \partial/\partial t \text{ with respect to } t, \\ \partial_{x_i}: & \text{ the usual derivation } \partial/\partial x_i \text{ with respect to } x_i, \\ \mathbf{x}: & \{x_1, \dots, x_n\} \\ \partial_{\mathbf{x}}: & \{\partial_{x_1}, \dots, \partial_{x_n}\}, \end{aligned}$$

The following formulas will also be frequently used:

$$\partial_x^\mu x^\nu = \begin{cases} \nu(\nu-1)\cdots(\nu-\mu+1)x^{\nu-\mu} + * \partial_x, & \nu \geq \mu \\ * \partial_x, & \nu < \mu \end{cases} \quad (1)$$

$$x^\mu \partial_x^\nu = \begin{cases} (-1)^\nu \mu(\mu-1)\cdots(\mu-\nu+1)x^{\mu-\nu} + \partial_x *, & \mu \geq \nu \\ \partial_x *, & \mu < \nu \end{cases} \quad (2)$$

where $* \in k\langle x, \partial_x \rangle$.

2 D -finite elements and differential forms

Throughout this paper, let k be an algebraically closed field of characteristic zero and let K be the differential field $k(t, x_1, \dots, x_n)$ with the derivations $\partial_t, \partial_{x_1}, \dots, \partial_{x_n}$. Let $\mathfrak{D} = K\langle \partial_t, \partial_{\mathbf{x}} \rangle$ be the ring of linear differential operators with coefficients in K . For $S \subset \{t, \mathbf{x}, \partial_t, \partial_{\mathbf{x}}\}$, denote by $k\langle S \rangle$ the subalgebra over k of \mathfrak{D} generated by S . For brevity, we denote $k\langle t, \mathbf{x}, \partial_t, \partial_{\mathbf{x}} \rangle$ by \mathfrak{W} . Let \mathcal{U} be the universal differential extension of K in which every algebraic differential equation having a solution in an extension of \mathcal{U} has a solution (see page 133 of [18] for more precise description).

Definition 2 *An element $f \in \mathcal{U}$ is said to be D -finite over K if for every $\delta \in \{\partial_t, \partial_{x_1}, \dots, \partial_{x_n}\}$, there is a nonzero operator $L_\delta \in K\langle \delta \rangle$ such that $L_\delta(f) = 0$.*

Denote by R the ring of D -finite elements over K , and by \mathcal{M} a free R -module of rank n with base $\{\mathbf{a}_1, \dots, \mathbf{a}_n\}$. Define a map $\mathfrak{D} \times \mathcal{M} \rightarrow \mathcal{M}$ given by

$$\left(L, \sum_{i=1}^n f_i \mathbf{a}_i \right) \rightarrow L \left(\sum_{i=1}^n f_i \mathbf{a}_i \right) := \sum_{i=1}^n L(f_i) \mathbf{a}_i.$$

This map endows \mathcal{M} with a left \mathfrak{D} -module structure. Let

$$\bigwedge(\mathcal{M}) = \bigoplus_{i=0}^n \bigwedge^i(\mathcal{M})$$

be the exterior algebra of \mathcal{M} , where $\bigwedge^i(\mathcal{M})$ denotes the i -th homogeneous part of $\bigwedge(\mathcal{M})$ as a graded R -algebra. We call an element in $\bigwedge^i(\mathcal{M})$ an i -form. Note that $\bigwedge(\mathcal{M})$ inherits a left \mathfrak{D} -module structure from \mathcal{M} . In fact, for $L \in \mathfrak{D}$ and $\omega = \sum f_{j_1, \dots, j_i} \mathbf{a}_{j_1} \wedge \cdots \wedge \mathbf{a}_{j_i} \in \bigwedge^i(\mathcal{M})$, one can define

$$L(\omega) = \sum L(f_{j_1, \dots, j_i}) \mathbf{a}_{j_1} \wedge \cdots \wedge \mathbf{a}_{j_i}$$

and for $\omega = \sum_i \omega_i$ with $\omega_i \in \wedge^i(\mathcal{M})$, define $L(\omega) = \sum_i L(\omega_i)$. Let $d : R \rightarrow \mathcal{M}$ be a map defined as

$$df = \partial_{x_1}(f)\mathbf{a}_1 + \cdots + \partial_{x_n}(f)\mathbf{a}_n$$

for any $f \in R$. Then d is a derivation over k . Note that for each $i = 1, \dots, n$, $dx_i = \mathbf{a}_i$. Hence in the rest of this paper we shall use $\{dx_1, \dots, dx_n\}$ instead of $\{\mathbf{a}_1, \dots, \mathbf{a}_n\}$. The map d can be extended to a derivation on $\wedge(\mathcal{M})$ which is defined recursively as

$$d(\omega_1 \wedge \omega_2) = d\omega_1 \wedge \omega_2 + (-1)^i \omega_1 \wedge d\omega_2$$

for any $\omega_1 \in \wedge^i(\mathcal{M})$ and $\omega_2 \in \wedge^j(\mathcal{M})$. For detailed definitions on exterior algebra and differential forms, we refer the readers to Chapter 19 of [22] and Chapter 1 of [35] respectively. As the usual differential forms, we introduce the following definition.

Definition 3 Let $\omega \in \wedge(\mathcal{M})$ be a form.

- (1) ω is said to be closed if $d\omega = 0$, and exact if there is $\eta \in \wedge(\mathcal{M})$ such that $\omega = d\eta$.
- (2) ω is said to be ∂_t -closed (resp. ∂_t -exact) if there is a nonzero $L \in k(t)\langle\partial_t\rangle$ such that $L(\omega)$ is closed (resp. exact).

Definition 4 Assume that $\omega \in \wedge(\mathcal{M})$. A nonzero $L \in k(t)\langle\partial_t\rangle$ is called a telescope for ω if $L(\omega)$ is exact.

3 Parametrized Poincaré lemma

The famous Poincaré lemma states that if B is an open ball in \mathbb{R}^n , any smooth closed i -form ω defined on B is exact, for any integer i with $1 \leq i \leq n$. In this section, we shall prove the following lemma which can be viewed as a parametrized analogue of Poincaré lemma for $\wedge(\mathcal{M})$.

Lemma 5 (Parameterized Poincaré lemma) Let $\omega \in \wedge^p(\mathcal{M})$. If ω is ∂_t -closed then it is ∂_t -exact.

To prove the above lemma, we need some lemmas.

Lemma 6 (Lipshitz's lemma (Lemma 3 of [25])) Assume that f is a D -finite element over $k(\mathbf{x})$. For each $i = 1, 3, 4, \dots, n$, there is a nonzero operator $L \in k(x_1, x_3, \dots, x_n)\langle\partial_{x_2}, \partial_{x_i}\rangle$ such that $L(f) = 0$.

The following lemma is a generalization of Lipshitz's lemma.

Lemma 7 *Assume that f_1, \dots, f_m are D -finite elements over $k(\mathbf{x}, t)$ and*

$$S \subset \{t, x_1, \dots, x_n, \partial_t, \partial_{x_1}, \dots, \partial_{x_n}\}$$

with $|S| > n + 1$. Then one can compute a nonzero operator L in $k\langle S \rangle$ such that $L(f_i) = 0$ for all $i = 1, \dots, m$.

Proof. For each $\delta \in \{\partial_t, \partial_{x_1}, \dots, \partial_{x_n}\}$ and $i = 1, \dots, m$, let $T_{i,\delta}$ be a nonzero operator in $K\langle \delta \rangle$ such that $T_{i,\delta}(f_i) = 0$. Set T_δ to be the least common left multiple of $T_{1,\delta}, \dots, T_{m,\delta}$. Then $T_\delta(f_i) = 0$ for all $i = 1, \dots, m$ and $\delta \in \{\partial_t, \partial_{x_1}, \dots, \partial_{x_n}\}$. The lemma then follows from an argument similar to that in the proof of Lipshitz's lemma. \blacksquare

Remark 8 *Lemma 7 originally appears in [39] (see Lemma 4.1), where Zeilberger proves the existence of the operator L in the setting of Weyl algebra and gives an algorithm to compute L in the case of two variables. Furthermore, there is a Mathematica package called *HolonomicFunctions* developed by Christoph which allows one to compute L (see [19]).*

Lemma 9 *Assume that f_1, \dots, f_m are D -finite over $k(\mathbf{x}, t)$, $I, J \subset \{1, \dots, n\}$ and $I \cap J = \emptyset$. Assume further that $V \subset \{x_i, \partial_{x_i} \mid i \in \{1, \dots, n\} \setminus (I \cup J)\}$ with $|V| = n - |I| - |J|$. Then one can compute an operator P of the form*

$$L + \sum_{i \in I} \partial_{x_i} M_i + \sum_{j \in J} N_j \partial_{x_j}$$

such that $P(f_l) = 0$ for all $l = 1, \dots, m$, where L is a nonzero operator in $k\langle \{t, \partial_t\} \cup V \rangle$, $M_i, N_j \in \mathfrak{W}$ and N_j is free of x_i for all $i \in I$ and $j \in J$.

Proof. Without loss of generality, we assume that $I = \{1, \dots, r\}$ and $J = \{r + 1, \dots, r + s\}$ where $r = |I|$ and $s = |J|$. Let

$$S = \{t, \partial_t\} \cup \{\partial_{x_i} \mid i \in I\} \cup \{x_j \mid j = r + 1, \dots, r + s\} \cup V.$$

Then $|S| = n + 2 > n + 1$. By Lemma 7, one can compute a $T \in k\langle S \rangle \setminus \{0\}$ such that $T(f_l) = 0$ for all $l = 1, \dots, m$. Write

$$T = \sum_{\mathbf{d}=(d_1, \dots, d_r) \in \Gamma_1} \partial_{x_1}^{d_1} \dots \partial_{x_r}^{d_r} T_{\mathbf{d}}$$

where $T_{\mathbf{d}} \in k\langle \{t, \partial_t, x_{r+1}, \dots, x_{r+s}\} \cup V \rangle \setminus \{0\}$ and Γ_1 is a finite subset of \mathbb{Z}^r . Let $\bar{\mathbf{d}} = (\bar{d}_1, \dots, \bar{d}_r)$ be the minimal element of Γ_1 with respect to the

lex order on \mathbb{Z}^r . Multiplying T by $\prod_{i=1}^r x_i^{\bar{d}_i}$ on the left and using the formula (2) yield that

$$\left(\prod_{i=1}^r x_i^{\bar{d}_i} \right) T = \alpha T_{\bar{\mathbf{d}}} + \sum_{i=1}^r \partial_{x_i} \tilde{T}_i \quad (3)$$

where α is a nonzero integer and $\tilde{T}_i \in k\langle S \cup \{x_i | i \in I\} \rangle$. Write

$$T_{\bar{\mathbf{d}}} = \sum_{\mathbf{e}=(e_1, \dots, e_s) \in \Gamma_2} L_{\mathbf{e}} x_{r+1}^{e_1} \dots x_{r+s}^{e_s}$$

where $L_{\mathbf{e}} \in k\langle \{t, \partial_t\} \cup V \rangle \setminus \{0\}$ and Γ_2 is a finite subset of \mathbb{Z}^s . Let $\bar{\mathbf{e}} = (\bar{e}_1, \dots, \bar{e}_s)$ be the maximal element of Γ_2 with respect to the lex order on \mathbb{Z}^s . Multiplying $T_{\bar{\mathbf{d}}}$ by $\prod_{i=1}^s \partial_{x_{r+i}}^{\bar{e}_i}$ on the left and using the formula (1) yield that

$$\left(\prod_{i=1}^s \partial_{x_{r+i}}^{\bar{e}_i} \right) T_{\bar{\mathbf{d}}} = \beta L_{\bar{\mathbf{e}}} + \sum_{j \in J} \tilde{L}_j \partial_{x_j} \quad (4)$$

where $\tilde{L}_i \in k\langle \{t, \partial_t, x_{r+1}, \dots, x_{r+s}, \partial_{x_{r+1}}, \dots, \partial_{x_{r+s}}\} \cup V \rangle$ and β is a nonzero integer. Combining (3) with (4) yields the required operator P . \blacksquare

Corollary 10 *Assume that f_1, \dots, f_m are D -finite over $k(\mathbf{x}, t)$, J is a subset of $\{1, \dots, n\}$ and $V \subset \{x_i, \partial_{x_i} | i \in \{1, \dots, n\} \setminus J\}$ with $|V| = n - |J|$. Assume further that $\partial_{x_j}(f_l) = 0$ for all $j \in J$ and $l = 1, \dots, m$. Then one can compute a nonzero $L \in k\langle \{t, \partial_t\} \cup V \rangle$ such that $L(f_l) = 0$ for all $l = 1, \dots, m$.*

Proof. In Lemma 9, set $I = \emptyset$. \blacksquare

The main result of this section is the following theorem which can be viewed as a generalization of Corollary 10 to differential forms. To describe and prove this theorem, let us recall some notation from the first chapter of [35]. For any $f \in R$, we define $d_0(f) = 0$ and

$$d_s(f) = \partial_{x_1}(f) dx_1 + \dots + \partial_{x_s}(f) dx_s$$

for $s \in \{1, 2, \dots, n\}$. We can extend d_s to the module $\wedge(\mathcal{M})$ in a natural way. Precisely, let $\omega = \sum_{i=1}^m f_i \mathbf{m}_i$ where \mathbf{m}_i is a monomial in dx_1, \dots, dx_n . Then $d_0(\omega) = 0$ and

$$d_s(\omega) = \sum_{i=1}^m \sum_{j=1}^s \partial_{x_j}(f_i) dx_j \wedge \mathbf{m}_i = \sum_{j=1}^s dx_j \wedge \partial_{x_j}(\omega).$$

By definition, one sees that

$$d_s(u \wedge dx_s) = d_{s-1}(u) \wedge dx_s \text{ and } d_s(u) = d_{s-1}(u) + dx_s \wedge \partial_{x_s}(u).$$

Theorem 11 Assume that $0 \leq s \leq n$, $V \subset \{x_{s+1}, \dots, x_n, \partial_{x_{s+1}}, \dots, \partial_{x_n}\}$ with $|V| = n - s$ and $\omega \in \wedge^p(\mathcal{M})$. If $d_s \omega = 0$, then one can compute a nonzero $L \in k\langle\{t, \partial_t\} \cup V\rangle$ and $\mu \in \wedge^{p-1}(\mathcal{M})$ such that $L(\omega) = d_s \mu$.

Remark 12 1. If $p = 0$, then $\omega = f \in R$ and $d_s f = 0$ if and only if $s = 0$ or $\partial_{x_i}(f) = 0$ for all $1 \leq i \leq s$ if $s > 0$. Therefore Corollary 10 is a special case of Theorem 11.

2. Note that the parametrized Poincaré lemma is just the special case of Theorem 11 when $s = n$.

Proof. We proceed by induction on s . Assume that $s = 0$ and write

$$\omega = \sum_{i=1}^m f_i \mathbf{m}_i$$

where \mathbf{m}_i a monomial in dx_1, dx_2, \dots, dx_n and $f_i \in R$. By Corollary 10 with $I = \emptyset$, one can compute a nonzero $L \in k\langle\{t, \partial_t\} \cup V\rangle$ such that $L(f_i) = 0$ for all $i = 1, \dots, m$. Then one has that

$$L(\omega) = \sum_{i=1}^m L(f_i) \mathbf{m}_i = 0.$$

This proves the base case. Now assume that the theorem holds for $s < \ell$ and consider the case $s = \ell$. Write

$$\omega = u \wedge dx_\ell + v$$

where both u and v do not involve dx_ℓ . Then the assumption $d_\ell \omega = 0$ implies that

$$d_{\ell-1} u \wedge dx_\ell + d_\ell v = d_{\ell-1} u \wedge dx_\ell + d_{\ell-1} v + dx_\ell \wedge \partial_{x_\ell}(v) = 0.$$

Since all of $d_{\ell-1} u, d_{\ell-1} v, \partial_{x_\ell}(v)$ do not involve dx_ℓ , one has that $d_{\ell-1} v = 0$ and $d_{\ell-1}(u) - \partial_{x_\ell}(v) = 0$. By the induction hypothesis, one can compute a nonzero $\tilde{L} \in k\langle\{t, x_\ell, \partial_t\} \cup V\rangle$ and $\tilde{\mu} \in \wedge^{p-1}(\mathcal{M})$ such that

$$\tilde{L}(v) = d_{\ell-1}(\tilde{\mu}). \quad (5)$$

We claim that \tilde{L} can be chosen to be free of x_ℓ . Write

$$\tilde{L} = \sum_{j=0}^d N_j x_\ell^j$$

where $N_j \in k\langle\{t, \partial_t\} \cup V\rangle$ and $N_d \neq 0$. Multiplying \tilde{L} by $\partial_{x_\ell}^d$ on the left and using the formula (2) yield that

$$\partial_{x_\ell}^d \tilde{L} = \sum_{j=0}^d N_j \partial_{x_\ell}^d x_\ell^j = \alpha N_d + \tilde{N} \partial_{x_\ell} \quad (6)$$

where α is a nonzero integer and $\tilde{N} \in k\langle\{t, x_\ell, \partial_t, \partial_{x_\ell}\} \cup V\rangle$. The equalities (5) and (6) together with $\partial_{x_\ell}(v) = d_{\ell-1}(\tilde{u})$ yield that $N_d(v) = d_{\ell-1}(\pi)$ for some $\pi \in \Lambda^{p-1}(\mathcal{M})$. This proves the claim. Now one has that

$$\tilde{L}(\omega) = \tilde{L}(u) \wedge dx_\ell + d_{\ell-1}(\tilde{\mu}) = \tilde{L}(u) \wedge dx_\ell + dx_\ell \wedge \partial_{x_\ell}(\tilde{\mu}) + d_\ell(\tilde{\mu}).$$

Since \tilde{L} is free of x_1, \dots, x_ℓ , $\tilde{L}d_\ell = d_\ell\tilde{L}$. This implies that

$$\begin{aligned} 0 &= \tilde{L}(d_\ell(\omega)) = d_\ell(\tilde{L}(\omega)) = d_{\ell-1}(\tilde{L}(u)) \wedge dx_\ell + dx_\ell \wedge d_{\ell-1}(\partial_{x_\ell}(\tilde{\mu})) \\ &= d_{\ell-1}(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) \wedge dx_\ell. \end{aligned}$$

Note that $\tilde{\mu}$ can always be chosen to be free of dx_ℓ . Hence one has that $d_{\ell-1}(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) = 0$. By the induction hypothesis, one can compute a nonzero $\bar{L} \in k\langle\{t, \partial_{x_\ell}, \partial_t\} \cup V\rangle$ and $\bar{\mu} \in \Lambda^{p-1}(\mathcal{M})$ such that

$$\bar{L}(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) = d_{\ell-1}(\bar{\mu}). \quad (7)$$

Write

$$\bar{L} = \sum_{j=e_1}^{e_2} \partial_{x_\ell}^j M_j$$

where $M_j \in k\langle\{t, \partial_t\} \cup V\rangle$ and $M_{e_1} \neq 0$. Multiplying \bar{L} by $x_\ell^{e_1}$ on the left and using the formula (2) yield that

$$x_\ell^{e_1} \bar{L} = \beta M_{e_1} + \partial_{x_\ell} \tilde{M}$$

where β is a nonzero integer and $\tilde{M} \in k\langle\{t, \partial_t, \partial_{x_\ell}, x_\ell\} \cup V\rangle$. Hence applying $x_\ell^{e_1}$ to the equality (7), one gets that

$$\beta M_{e_1}(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) = d_{\ell-1}(x_\ell^{e_1} \bar{\mu}) + \partial_{x_\ell}(\tilde{M}(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu}))).$$

Set $L = \beta M_{e_1} \tilde{L}$. Then one has that

$$\begin{aligned} L(\omega) &= \beta M_{e_1} \left((\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) \wedge dx_\ell + d_\ell(\tilde{\mu}) \right) \\ &= \left(\beta M_{e_1} (\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu})) \right) \wedge dx_\ell + d_\ell(\beta M_{e_1}(\tilde{\mu})) \\ &= d_{\ell-1}(x_\ell^{e_1} \bar{\mu}) \wedge dx_\ell + \partial_{x_\ell} \tilde{M} \left(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu}) \right) \wedge dx_\ell + d_\ell(\beta M_{e_1}(\tilde{\mu})) \\ &= d_\ell \left(x_\ell^{e_1} \bar{\mu} + \tilde{M} \left(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu}) \right) + \beta M_{e_1}(\tilde{\mu}) \right). \end{aligned}$$

The last equality holds because

$$d_{\ell-1} \left(\tilde{M} \left(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu}) \right) \right) = \tilde{M} d_{\ell-1} \left(\tilde{L}(u) - \partial_{x_\ell}(\tilde{\mu}) \right) = 0.$$

■

Remark 13 *Lemma 5 can be derived from the finiteness of the de Rham cohomology groups of D -modules in the Bernstein class. To see this, let ω be a differential s -form with coefficients in R and let M be the D -module generated by all coefficients of ω and all derivatives of these coefficients with respect to ∂_t . By Proposition 5.2 on page 12 of [5], M is a D -module in the Bernstein class. Assume that ω is closed. Then $\partial_t^j(\omega) \in H_{DR}^s(M)$, the s -th de Rham cohomology group of M , for all nonnegative integer j . By Theorem 6.1 on page 16 of [5], $H_{DR}^s(M)$ is of finite dimension over $k(t)$. This implies that there are $a_0, \dots, a_m \in k(t)$ such that $\sum_{j=0}^m a_j \partial_t^j(\omega) = 0$ in $H_{DR}^s(M)$, i.e. $\sum_{j=0}^m a_j \partial_t^j(\omega)$ is exact. This proves the existence of telescopers for the ∂_t -closed differential forms. However the proof of Theorem 11 is constructive and it provides a method to compute a telescoper if it exists.*

The proof of Theorem 11 can be summarized as the following algorithm.

Algorithm 14 *Input: $\omega \in \Lambda^p(\mathcal{M})$ and $V \subset \{x_i, \partial_{x_i} | i = s+1, \dots, n\}$ satisfying that $d_s(\omega) = 0$ and $|V| = n - s$
Output: a nonzero $L \in k\langle\{t, \partial_t\} \cup V\rangle$ such that $L(\omega) = d_s(\mu)$.*

1. If $\omega \in R$, then by Corollary 10, compute a nonzero $L \in k\langle\{t, \partial_t\} \cup V\rangle$ such that $L(\omega) = 0$. Return L .
2. Write $\omega = u \wedge dx_s + v$ with u, v not involving dx_s .
3. Call Algorithm 14 with v and $V \cup \{x_s\}$ as inputs and let \tilde{L} be the output.
 - (a) Write $\tilde{L} = \sum_{j=0}^d N_j x_s^j$ with $N_j \in k\langle\{t, \partial_t\} \cup V\rangle$ and $N_d \neq 0$.
 - (b) Compute a $\tilde{\mu} \in \Lambda^{p-1}(\mathcal{M})$ such that $N_d(v) = d_{s-1}(\tilde{\mu})$.
4. Write $N_d(\omega) = (N_d(u) - \partial_{x_s}(\tilde{\mu})) \wedge dx_s + d_s(\tilde{\mu})$.
5. Call Algorithm 14 with $N_d(u) - \partial_{x_s}(\tilde{\mu})$ and $V \cup \{\partial_{x_s}\}$ as inputs and let \bar{L} be the output.
6. Write $\bar{L} = \sum_{j=e_1}^{e_2} \partial_{x_s}^j M_j$ with $M_j \in k\langle\{t, \partial_t\} \cup V\rangle$ and $M_{e_1} \neq 0$.
7. Return $M_{e_1} N_d$.

4 The existence of telescopers

It is easy to see that if a differential form is ∂_t -exact then it is ∂_t -closed. Therefore Lemma 5 implies that given a $\omega \in \bigwedge^p(\mathcal{M})$, to decide whether it has a telescoper, it suffices to decide whether there is a nonzero $L \in k\langle t, \partial_t \rangle$ such that $L(d\omega) = 0$. Suppose that

$$d\omega = \sum_{1 \leq i_1 < \dots < i_{p+1} \leq n} a_{i_1, \dots, i_{p+1}} dx_{i_1} \cdots dx_{i_{p+1}}, \quad a_{i_1, \dots, i_{p+1}} \in \mathcal{U}.$$

Then $L(d\omega) = 0$ if and only if $L(a_{i_1, \dots, i_{p+1}}) = 0$ for all $1 \leq i_1 < \dots < i_{p+1} \leq n$. So the existence problem of telescopers can be reduced to the following problem.

Problem 15 *Given an element $f \in R$ and its minimal annihilating operator $P \in K\langle \partial_t \rangle$, decide whether there exists a nonzero $L \in k\langle t, \partial_t \rangle$ such that $L(f) = 0$.*

Note that f is annihilated by a nonzero $L \in k(t)\langle \partial_t \rangle$ if and only if P is a right-hand factor of L , i.e. $L = QP$ for some $Q \in K\langle \partial_t \rangle$. For convenience, we introduce the following definition.

Definition 16 *An operator $P \in K\langle \partial_t \rangle$ is called (\mathbf{x}, t) -separable if there is a nonzero $L \in k(t)\langle \partial_t \rangle$ such that $L = QP$ for some $Q \in K\langle \partial_t \rangle$.*

Problem 15 then is reduced to the following one.

Problem 17 *Given a $P \in K\langle \partial_t \rangle \setminus \{0\}$, decide whether P is (\mathbf{x}, t) -separable.*

The rest of this paper is aimed at developing an algorithm to solve the above problem. Let us first investigate the solutions of (\mathbf{x}, t) -separable operators.

Notation 18

$$C_t := \{c \in \mathcal{U} \mid \partial_t(c) = 0\}, \quad C_{\mathbf{x}} := \{c \in \mathcal{U} \mid \forall x \in \mathbf{x}, \partial_x(c) = 0\}.$$

Assume that $L \in k(t)\langle \partial_t \rangle \setminus \{0\}$. By Corollary 1.2.12 of [30], the solution space of $L = 0$ in \mathcal{U} is a C_t -vector space of dimension $\text{ord}(L)$. Moreover we have the following lemma.

Lemma 19 *If $L \in k(t)\langle \partial_t \rangle \setminus \{0\}$, then the solution space of $L = 0$ in \mathcal{U} has a basis in $C_{\mathbf{x}}$.*

Proof. Let A_0 be the companion matrix of $L(y) = 0$, i.e.

$$A_0 = \begin{pmatrix} 0 & 1 & & & \\ & 0 & 1 & & \\ & & \ddots & \ddots & \\ & & & 0 & 1 \\ -a_0 & -a_1 & -a_2 & \dots & -a_{m-1} \end{pmatrix}$$

where $m = \text{ord}(L)$ and $L = \partial_t^m + a_{m-1}\partial_t^{m-1} + \dots + a_0$. Set $A_i = 0$ for all $i = 1, \dots, n$. Let $\partial_0 = \partial_t, \partial_i = \partial_{x_i}$ for $i = 1, \dots, n$. Then the system

$$\partial_0(Y) = A_0Y, \partial_1(Y) = A_1Y, \dots, \partial_n(Y) = A_nY$$

satisfies the integrability condition:

$$\partial_i(A_j) - \partial_j(A_i) = A_iA_j - A_jA_i$$

for all $0 \leq i < j \leq n$. Therefore there is a solution V in $\text{GL}_m(\mathcal{U})$. Let \mathbf{v} be the first row of V . Note that $\det(V)$ is the Wronskian determinant of \mathbf{v} and $\det(V) \neq 0$. These imply that \mathbf{v} is a basis of the solution space of $L(y) = 0$. Since $\partial_i(\mathbf{v}) = 0$ for all $1 \leq i \leq n$, \mathbf{v} has entries in $C_{\mathbf{x}}$. \blacksquare

As a consequence, we have the following corollary.

Corollary 20 *Assume that $P \in K\langle\partial_t\rangle \setminus \{0\}$. Then P is (\mathbf{x}, t) -separable if and only if the solutions of $P(y) = 0$ in \mathcal{U} are of the form*

$$\sum_{i=1}^s g_i h_i, \quad g_i \in C_t, h_i \in C_{\mathbf{x}} \cap \{f \in \mathcal{U} \mid Q(f) = 0\} \quad (8)$$

for some $Q \in K\langle\partial_t\rangle$.

Proof. The ‘‘only if’’ part is a direct consequence of Lemma 19. For the ‘‘if’’ part, one only need to prove that if $h \in C_{\mathbf{x}} \cap \{f \in \mathcal{U} \mid Q(f) = 0\}$ then h is annihilated by a nonzero operator in $k(t)\langle\partial_t\rangle$. Suppose that $h \in C_{\mathbf{x}} \cap \{f \in \mathcal{U} \mid Q(f) = 0\}$. Let L be the monic operator in $K\langle\partial_t\rangle \setminus \{0\}$ which annihilates h and is of minimal order. Write

$$L = \partial_t^\ell + \sum_{i=0}^{\ell-1} a_i \partial_t^i, \quad a_i \in K.$$

Then for every $j \in \{1, \dots, n\}$

$$0 = \partial_{x_j}(L(h)) = \sum_{i=0}^{\ell-1} \partial_{x_j}(a_i) \partial_t^i(h) + L(\partial_{x_j}(h)) = \sum_{i=0}^{\ell-1} \partial_{x_j}(a_i) \partial_t^i(h).$$

The last equality holds because $h \in C_{\mathbf{x}}$. By the minimality of L , one sees that $\partial_{x_j}(a_i) = 0$ for all $i = 0, \dots, \ell - 1$ and all $j = 1, \dots, n$. Hence $a_i \in k(t)$ for all i . In other words, $L \in k(t)\langle \partial_t \rangle$. \blacksquare

For convention, we introduce the following definition.

- Definition 21** (1) We say $f \in \mathcal{U}$ is split if it can be written as the form $f = gh$ where $g \in C_t$ and $h \in C_{\mathbf{x}}$, and say f is semisplit if it is the sum of finitely many split elements.
- (2) We say a nonzero operator $P \in K\langle \partial_t \rangle$ is semisplit if it is monic and all its coefficients are semisplit.

The semisplit operators have the following property.

Lemma 22 Assume that $P = Q_1 Q_2$ where P, Q_1, Q_2 are monic operators in $K\langle \partial_t \rangle$. Assume further that $Q_2 \in k(t)[\mathbf{x}, 1/r]\langle \partial_t \rangle$ where $r \in k[\mathbf{x}, t]$. Then $P \in k(t)[\mathbf{x}, 1/r]\langle \partial_t \rangle$ if and only if Q_1 is also.

Proof. Comparing the coefficients on both sides of $P = Q_1 Q_2$ concludes the lemma. \blacksquare

As a direct consequence, we have the following corollary.

Corollary 23 Assume that $P = Q_1 Q_2$ where P, Q_1, Q_2 are monic operators in $K\langle \partial_t \rangle$. Assume further that Q_2 is semisplit. Then P is semisplit if and only if Q_1 is also.

4.1 The completely reducible case

In Proposition 10 of [11], we show that given a hyperexponential function h over K , $\text{ann}(h) \cap k(t)\langle \partial_t \rangle \neq \{0\}$ if and only if there is a nonzero $p \in k(\mathbf{x})[t]$ and $r \in k(t)$ such that

$$a = \frac{\partial_t(p)}{p} + r,$$

where $a = \partial_t(h)/h$. Remark that a, p, r with $p \neq 0$ satisfy the above equality if and only if $\frac{1}{p}(\partial_t - a) = (\partial_t - r)\frac{1}{p}$. Under the notion of (\mathbf{x}, t) -separable

and the language of differential operators, Proposition 10 of [11] states that $\partial_t - a$ is (\mathbf{x}, t) -separable if and only if it is similar to a first order operator in $k(t)\langle\partial_t\rangle$ by some $1/p$ with p being nonzero polynomial in t . In this section, we shall generalize Proposition 10 of [11] to the case of completely reducible operators. We shall use $\text{lcm}(Q_1, Q_2)$ to denote the monic operator of minimal order which is divisible by both Q_1 and Q_2 on the right. We shall prove that if P is (\mathbf{x}, t) -separable and completely reducible then there is a nonzero $L \in k(t)\langle\partial_t\rangle$ such that P is the transformation of L by some Q with semisplit coefficients. To this end, we need to introduce some notations from [28].

Definition 24 *Assume that $P, Q \in K\langle\partial_t\rangle \setminus \{0\}$.*

1. *We say \tilde{P} is the transformation of P by Q if \tilde{P} is the monic operator satisfying that $\tilde{P}Q = \lambda\text{lcm}(P, Q)$ for some $\lambda \in K$.*
2. *We say \tilde{P} is similar to P (by Q) if there is an operator Q with $\text{gcd}(P, Q) = 1$ such that \tilde{P} is the transformation of P by Q , where $\text{gcd}(P, Q)$ denotes the greatest common right-hand factor of P and Q .*

Definition 25 1. *We say $P \in K\langle\partial_t\rangle$ is completely reducible if it is the lcm of a family of irreducible operators in $K\langle\partial_t\rangle$.*

2. *We say $Q \in K\langle\partial_t\rangle$ is the maximal completely reducible right-hand factor of $P \in K\langle\partial_t\rangle$ if Q is the lcm of all irreducible right-hand factors of P .*

Given a $P \in K\langle\partial_t\rangle$, Theorem 7 of [28] implies that P has the following unique decomposition called the maximal completely reducible decomposition or the m.c.r. decomposition for short,

$$P = \lambda H_r H_{r-1} \dots H_1$$

where $\lambda \in K$ and H_i is the maximal completely reducible right-hand factor of $H_r \dots H_i$. For an $L \in k(t)\langle\partial_t\rangle$, it has two m.c.r. decompositions viewed it as an operator in $k(t)\langle\partial_t\rangle$ and an operator in $K\langle\partial_t\rangle$ respectively. In the following, we shall prove that these two decompositions coincide. For convenience, we shall denote by $P_{x_i=c_i}$ the operator obtained by replacing x_i by $c_i \in k$ in P .

Lemma 26 *Assume that P, L are two monic operators in $K\langle\partial_t\rangle$. Assume further that $P \in k(t)[\mathbf{x}, 1/r]\langle\partial_t\rangle$ with $r \in k[\mathbf{x}, t]$, and $L \in k(t)\langle\partial_t\rangle$. Let $\mathbf{c} \in k^n$ be such that $r(\mathbf{c}) \neq 0$.*

1. If $\text{gcd}(P_{\mathbf{x}=\mathbf{c}}, L) = 1$ then $\text{gcd}(P, L) = 1$.
2. If $\text{gcd}(P, L) = 1$ then there is $\mathbf{a} \in k^n$ such that $r(\mathbf{a}) \neq 0$ and $\text{gcd}(P_{\mathbf{x}=\mathbf{a}}, L) = 1$.

Proof. 1. We shall prove the lemma by induction on $n = |\mathbf{x}|$. Assume that $n = 1$, and $\text{gcd}(P, L) \neq 1$. Then there are $M, N \in k(t)[x_1]\langle \partial_t \rangle$ with $\text{ord}(M) < \text{ord}(L)$ such that $MP + NL = 0$. Write

$$M = \sum_{i=0}^{m-1} a_i \partial_t^i, \quad N = \sum_{i=0}^s b_i \partial_t^i$$

where $m = \text{ord}(L)$. If the a_i 's have a common factor c in $k(t_1)[x_1]$, then one sees that c is a common factor of the b_i 's. Thus we can cancel this factor c . So without loss of generality, we may assume that the a_i 's have no common factor. This implies that $M_{x_1=c_1} \neq 0$ and $M_{x_1=c_1}P_{x_1=c_1} + N_{x_1=c_1}L = 0$. Since $\text{ord}(M_{x_1=c_1}) < \text{ord}(L)$, $\text{gcd}(P_{x_1=c_1}, L) \neq 1$, a contradiction. For the general case, set $Q = P_{x_1=c_1}$. Then $Q_{x_2=c_2, \dots, x_n=c_n} = P_{\mathbf{x}=\mathbf{c}}$. This implies that $\text{gcd}(Q_{x_2=c_2, \dots, x_n=c_n}, L) = 1$. By the induction hypothesis, $\text{gcd}(Q, L) = 1$. Finally, regarding P and L as operators with coefficients in $k(t, x_2, \dots, x_n)[x_1, 1/r]$ and by the induction hypothesis again, we get $\text{gcd}(P, L) = 1$.

2. Since $\text{gcd}(P, L) = 1$, there are $M, N \in K\langle \partial_t \rangle$ such that $MP + NL = 1$. Let $\mathbf{a} \in k^n$ be such that $r(\mathbf{a}) \neq 0$ and both $M_{\mathbf{x}=\mathbf{a}}$ and $N_{\mathbf{x}=\mathbf{a}}$ are well-defined. For such \mathbf{a} , one has that $M_{\mathbf{x}=\mathbf{a}}P_{\mathbf{x}=\mathbf{a}} + N_{\mathbf{x}=\mathbf{a}}L = 1$ and then $\text{gcd}(P_{\mathbf{x}=\mathbf{a}}, L) = 1$. ■

Lemma 27 *Let $L \in k(t)\langle \partial_t \rangle$. The m.c.r. decompositions of L viewed as an operator in $k(t)\langle \partial_t \rangle$ and an operator in $K\langle \partial_t \rangle$ respectively coincide.*

Proof. We first claim that an irreducible operator of $k(t)\langle \partial_t \rangle$ is irreducible in $K\langle \partial_t \rangle$. Let P be a monic irreducible operator in $k(t)\langle \partial_t \rangle$ and assume that Q is a monic right-hand factor of P in $K\langle \partial_t \rangle$ with $1 \leq \text{ord}(Q) < \text{ord}(P)$. Then $P = \tilde{Q}Q$ for some $\tilde{Q} \in K\langle \partial_t \rangle$. Suppose that $Q \in k(t)[\mathbf{x}, 1/r]\langle \partial_t \rangle$. By Lemma 22, \tilde{Q} belongs to $k(t)[\mathbf{x}, 1/r]\langle \partial_t \rangle$. Let $\mathbf{c} \in k^n$ be such that $r(\mathbf{c}) \neq 0$. Then $P = \tilde{Q}_{\mathbf{x}=\mathbf{c}}Q_{\mathbf{x}=\mathbf{c}}$ and $1 \leq \text{ord}(Q_{\mathbf{x}=\mathbf{c}}) < \text{ord}(P)$. These imply that P is reducible in $k(t)\langle \partial_t \rangle$, a contradiction. So P is irreducible in $K\langle \partial_t \rangle$ and thus the claim holds. Let $L = \lambda H_r H_{r-1} \dots H_1$ be the m.c.r. decomposition in $k(t)\langle \partial_t \rangle$. The above claim implies that H_1 viewed as an operator in $K\langle \partial_t \rangle$ is completely reducible. Assume that H_1 is not the maximal completely reducible right-hand factor of L in $K\langle \partial_t \rangle$. Let $M \in K\langle \partial_t \rangle \setminus K$ be a monic

irreducible right-hand factor of L satisfying that $\gcd(M, H_1) = 1$. Due to Lemma 26, there is $\mathbf{a} \in k^n$ satisfying that $\gcd(M_{\mathbf{x}=\mathbf{a}}, H_1) = 1$. Note that $M_{\mathbf{x}=\mathbf{a}}$ is a right-hand factor of L . Therefore $M_{\mathbf{x}=\mathbf{a}}$ has some irreducible right-hand factor of L as a right-hand factor. Such irreducible factor must be a right-hand factor of H_1 and thus $\gcd(M_{\mathbf{x}=\mathbf{a}}, H_1) \neq 1$, a contradiction. Therefore H_1 is the maximal completely reducible right-hand factor of L in $K\langle\partial_t\rangle$. Using the induction on the order, one sees that $\lambda H_r H_{r-1} \dots H_1$ is the m.c.r. decomposition of L in $K\langle\partial_t\rangle$. \blacksquare

Lemma 28 *Assume that P is monic, (\mathbf{x}, t) -separable and completely reducible. Assume further that $P \in k(t)[\mathbf{x}, 1/r]\langle\partial_t\rangle$ with $r \in k[\mathbf{x}, t]$. Let $\mathbf{c} \in k^n$ be such that $r(\mathbf{c}) \neq 0$. Then $P_{\mathbf{x}=\mathbf{c}}$ is similar to P .*

Proof. Let \tilde{L} be a nonzero monic operator in $k(t)\langle\partial_t\rangle$ with P as a right-hand factor. Since P is completely reducible, by Theorem 8 of [28], P is a right-hand factor of the maximal completely reducible right-hand factor of \tilde{L} . By Lemma 27, the maximal completely reducible right-hand factor of \tilde{L} is in $k(t)\langle\partial_t\rangle$. Hence we may assume that \tilde{L} is completely reducible after replacing \tilde{L} by its maximal completely reducible right-hand factor. Assume that $\tilde{L} = QP$ for some $Q \in K\langle\partial_t\rangle$. By Lemma 22, $Q \in k(t)[\mathbf{x}, 1/r]\langle\partial_t\rangle$. Then $\tilde{L} = Q_{\mathbf{x}=\mathbf{c}}P_{\mathbf{x}=\mathbf{c}}$, i.e. $P_{\mathbf{x}=\mathbf{c}}$ is a right-hand factor of \tilde{L} . We claim that for a right-hand factor T of \tilde{L} , there is a right-hand factor L of \tilde{L} satisfying that $\gcd(T, L) = 1$ and $\text{lcm}(T, L) = \tilde{L}$. We prove this claim by induction on $s = \text{ord}(\tilde{L}) - \text{ord}(T)$. When $s = 0$, there is nothing to prove. Assume that $s > 0$. Then since \tilde{L} is completely reducible, there is an irreducible right-hand factor L_1 of \tilde{L} such that $\gcd(T, L_1) = 1$. Let $N = \text{lcm}(T, L_1)$. We have that $\text{ord}(N) = \text{ord}(T) + \text{ord}(L_1)$. Therefore $\text{ord}(\tilde{L}) - \text{ord}(N) < s$. By induction hypothesis, there is a right-hand factor L_2 of \tilde{L} such that $\gcd(N, L_2) = 1$ and $\text{lcm}(N, L_2) = \tilde{L}$. Let $L = \text{lcm}(L_1, L_2)$. Then

$$\tilde{L} = \text{lcm}(N, L_2) = \text{lcm}(T, L_1, L_2) = \text{lcm}(T, L).$$

Taking the order of the operators in the above equality yields that

$$\begin{aligned} \text{ord}(\text{lcm}(T, L)) &= \text{ord}(\text{lcm}(N, L_2)) = \text{ord}(N) + \text{ord}(L_2) \\ &= \text{ord}(T) + \text{ord}(L_1) + \text{ord}(L_2). \end{aligned}$$

On the other hand, we have

$$\text{ord}(\text{lcm}(T, L)) \leq \text{ord}(T) + \text{ord}(L) \leq \text{ord}(T) + \text{ord}(L_1) + \text{ord}(L_2).$$

This implies that

$$\text{ord}(\text{lclm}(T, L)) = \text{ord}(T) + \text{ord}(L).$$

So $\text{gcd}(T, L) = 1$ and then L is a required operator. This proves the claim. Now let $L_{\mathbf{c}}$ be a right-hand factor of \tilde{L} satisfying that $\text{gcd}(P_{\mathbf{x}=\mathbf{c}}, L_{\mathbf{c}}) = 1$ and $\text{lclm}(P_{\mathbf{x}=\mathbf{c}}, L_{\mathbf{c}}) = \tilde{L}$. Let $M \in k(t)\langle \partial_t \rangle$ be such that $\tilde{L} = ML_{\mathbf{c}}$. Then $P_{\mathbf{x}=\mathbf{c}}$ is similar to M . It remains to show that P is also similar to M . Due to Lemma 26, $\text{gcd}(P, L_{\mathbf{c}}) = 1$. Then

$$\text{ord}(\text{lclm}(P, L_{\mathbf{c}})) = \text{ord}(P) + \text{ord}(L_{\mathbf{c}}) = \text{ord}(P_{\mathbf{x}=\mathbf{c}}) + \text{ord}(L_{\mathbf{c}}) = \text{ord}(\tilde{L}).$$

Note that $\text{lclm}(P, L_{\mathbf{c}})$ is a right-hand factor of \tilde{L} . Hence $\text{lclm}(P, L_{\mathbf{c}}) = \tilde{L}$ and thus P is similar to M . \blacksquare

For the general case, the above lemma is not true anymore as shown in the following example.

Example 29 Let $y = x_1 \log(t+1) + x_2 \log(t-1)$ and

$$P = \partial_t^2 + \frac{(t-1)^2 x_1 + (t+1)^2 x_2}{(t^2-1)((t-1)x_1 + (t+1)x_2)} \partial_t.$$

Then P is (x, t) -separable since $\{1, y\}$ is a basis of the solution space of $P = 0$ in \mathcal{U} . We claim that P is not similar to $P_{\mathbf{x}=\mathbf{c}}$ for any $\mathbf{c} \in k^2 \setminus \{(0, 0)\}$. Suppose on the contrary that P is similar to $P_{\mathbf{x}=\mathbf{c}}$ for some $\mathbf{c} = (c_1, c_2) \in k^2 \setminus \{(0, 0)\}$, i.e. there are $a, b \in k(\mathbf{x}, t)$, not all zero, such that $\text{gcd}(a\partial_t + b, P_{\mathbf{x}=\mathbf{c}}) = 1$ and P is the transformation of $P_{\mathbf{x}=\mathbf{c}}$ by $a\partial_t + b$. Denote $Q = a\partial_t + b$. As $\{1, y_{\mathbf{x}=\mathbf{c}}\}$ is a basis of the solution space of $P_{\mathbf{x}=\mathbf{c}}$, $\{Q(1), Q(y_{\mathbf{x}=\mathbf{c}})\}$ is a basis of the solution space of $P = 0$. In other words, there is $C \in \text{GL}_2(C_t)$ such that

$$\left(b, a \left(\frac{c_1}{t+1} + \frac{c_2}{t-1} \right) + by_{\mathbf{x}=\mathbf{c}} \right) = (1, y)C.$$

Note that $\log(t+1), \log(t-1), 1$ are linearly independent over $k(x_1, x_2, t)$. We have that $b \in C_t \setminus \{0\}$ and $bc_1 = \tilde{c}x_1, bc_2 = \tilde{c}x_2$ for some $\tilde{c} \in C_t$. This implies that $x_1/x_2 = c_1/c_2 \in k$, a contradiction.

When the given two operators are of length two, i.e. they are the products of two irreducible operators, a criterion for the similarity is presented in [24]. For the general case, suppose that P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by Q . Then the operator Q is a solution of the following mixed differential equation

$$Pz \equiv 0 \pmod{P_{\mathbf{x}=\mathbf{c}}}. \quad (9)$$

An algorithm for computing all solutions of the above mixed differential equation is developed in [33]. In the following, we shall show that if P is (\mathbf{x}, t) -separable then Q is an operator with semisplit coefficients. Note that Q can be chosen to be of order less than $\text{ord}(P_{\mathbf{x}=\mathbf{c}})$ and all solutions of the mixed differential equation with order less than $\text{ord}(P_{\mathbf{x}=\mathbf{c}})$ form a vector space over $k(\mathbf{x})$ of finite dimension. Furthermore Q induces an isomorphism from the solution space of $P_{\mathbf{x}=\mathbf{c}}(y) = 0$ to that of $P(y) = 0$.

Proposition 30 *Assume that P is monic and completely reducible. Assume further that $P \in k(t)[\mathbf{x}, 1/r]\langle \partial_t \rangle$ with $r \in k[\mathbf{x}, t]$. Let $\mathbf{c} \in k^n$ be such that $r(\mathbf{c}) \neq 0$. Then P is (\mathbf{x}, t) -separable if and only if P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by an operator Q with semisplit coefficients.*

Proof. Denote $m = \text{ord}(P_{\mathbf{x}=\mathbf{c}}) = \text{ord}(P)$. Assume that $\{\alpha_1, \dots, \alpha_m\}$ is a basis of the solution space of $P_{\mathbf{x}=\mathbf{c}}(y) = 0$ in $C_{\mathbf{x}}$ and P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by Q . Write $Q = \sum_{i=0}^{m-1} a_i \partial_t^i$ where $a_i \in K$. Then

$$(Q(\alpha_1), \dots, Q(\alpha_m)) = (a_0, \dots, a_{m-1}) \begin{pmatrix} \alpha_1 & \alpha_2 & \dots & \alpha_m \\ \alpha'_1 & \alpha'_2 & \dots & \alpha'_m \\ \vdots & \vdots & & \vdots \\ \alpha_1^{(m-1)} & \alpha_2^{(m-1)} & \dots & \alpha_m^{(m-1)} \end{pmatrix}$$

and $Q(\alpha_1), \dots, Q(\alpha_m)$ form a basis of the solution space of $P(y) = 0$.

Now suppose that P is (\mathbf{x}, t) -separable. Due to Lemma 28, P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by Q . By Corollary 20, the $Q(\alpha_i)$ are semisplit. The above equalities then imply that the a_i are semisplit. Conversely, assume that P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by Q and the a_i are semisplit. It is easy to see the $Q(\alpha_i)$ are semisplit. By Corollary 20 again, P is (\mathbf{x}, t) -separable. \blacksquare

Using the algorithm developed in [33], we can compute a basis of the solution space over $k(\mathbf{x})$ of the equation (9). It is clear that the solutions with semisplit entries form a subspace. We can compute a basis for this subspace as follows. Suppose that $\{Q_1, \dots, Q_\ell\}$ is a basis of the solution space of the equation (9) consisting of solutions with order less than $\text{ord}(P_{\mathbf{x}=\mathbf{c}})$. We may identify Q_i with a vector $\mathbf{g}_i \in K^m$ under the basis $1, \partial_t, \dots, \partial_t^{m-1}$. Let $q \in k(\mathbf{x})[t]$ be a common denominator of all entries of the \mathbf{g}_i . Write $\mathbf{g}_i = \mathbf{p}_i/q$ for each $i = 1, \dots, \ell$, where $\mathbf{p}_i \in k(\mathbf{x})[t]^m$. Write $q = q_1 q_2$ where q_2 is split but q_1 is not. Note that a rational function in t with coefficients in $k(\mathbf{x})$ is semisplit if and only if its denominator is split. For $c_1, \dots, c_\ell \in k(\mathbf{x})$, $\sum_{i=1}^\ell c_i \mathbf{g}_i$ is semisplit if and only if all entries of $\sum_{i=1}^\ell c_i \mathbf{p}_i$ are divided by

q_1 . For $i = 1, \dots, \ell$, let \mathbf{h}_i be the vector whose entries are the remainders of the corresponding entries of \mathbf{p}_i by q_1 . Then all entries of $\sum_{i=1}^{\ell} c_i \mathbf{p}_i$ are divided by q_1 if and only if $\sum_{i=1}^{\ell} c_i \mathbf{h}_i = 0$. Let $\mathbf{c}_1, \dots, \mathbf{c}_s$ be a basis of the solution space of $\sum_{i=1}^{\ell} z_i \mathbf{h}_i = 0$. Then $\{(Q_1, \dots, Q_{\ell})\mathbf{c}_i \mid i = 1, \dots, s\}$ is the required basis. Consequently, the required basis can be computed by solving the system of linear equations $\sum_{i=1}^{\ell} z_i \mathbf{h}_i = 0$.

In the following, for the sake of notations, we assume that $\{Q_1, \dots, Q_{\ell}\}$ is a basis of the solution space of the equation (9) consisting of solutions with semisplit coefficients. By Proposition 30 and the definition of similarity, P is (\mathbf{x}, t) -separable if and only if there is a nonzero \tilde{Q} in the space spanned by Q_1, \dots, Q_{ℓ} such that $\text{gcd}(P_{\mathbf{x}=\mathbf{c}}, \tilde{Q}) = 1$. Note that \tilde{Q} induces a homomorphism from the solutions space of $P_{\mathbf{x}=\mathbf{c}}(y) = 0$ to that of $P(y) = 0$. Moreover, one can easily see that $\text{gcd}(P_{\mathbf{x}=\mathbf{c}}, \tilde{Q}) = 1$ if and only if \tilde{Q} is an isomorphism i.e. $\tilde{Q}(\alpha_1), \dots, \tilde{Q}(\alpha_m)$ form a basis of the solution space of $P(y) = 0$ where $\{\alpha_1, \dots, \alpha_m\}$ is a basis of the solution space of $P_{\mathbf{x}=\mathbf{c}}(y) = 0$. Assume that $\tilde{Q} = \sum_{i=0}^{m-1} a_{0,i} \partial_t^i$ with $a_{0,i} \in K$. Using the relation $P_{\mathbf{x}=\mathbf{c}}(\alpha_j) = 0$ with $j = 1, \dots, m$, one has that for all $j = 1, \dots, m$

$$\tilde{Q}(\alpha_j)' = \left(\sum_{i=0}^{m-1} a_{0,i} \alpha_j^{(i)} \right)' = \sum_{i=0}^{m-1} a_{1,i} \alpha_j^{(i)}$$

for some $a_{1,i} \in K$. Repeating this process, we can compute $a_{l,i} \in K$ such that for all $j = 1, \dots, m$ and $l = 1, \dots, m-1$,

$$\tilde{Q}(\alpha_j)^{(l)} = \sum_{i=0}^{m-1} a_{l,i} \alpha_j^{(i)}.$$

Now suppose that $\tilde{Q} = \sum_{i=1}^{\ell} z_i Q_i$ with $z_i \in k(\mathbf{x})$. One sees that the $a_{l,i}$ are linear in z_1, \dots, z_{ℓ} . Set $A(\mathbf{z}) = (a_{i,j})_{0 \leq i, j \leq m-1}$ with $\mathbf{z} = (z_1, \dots, z_{\ell})$. Then one has that

$$A(\mathbf{z}) \begin{pmatrix} \alpha_1 & \dots & \alpha_m \\ \vdots & & \vdots \\ \alpha^{(m-1)} & \dots & \alpha_m^{(m-1)} \end{pmatrix} = \begin{pmatrix} \tilde{Q}(\alpha_1) & \dots & \tilde{Q}(\alpha_m) \\ \vdots & & \vdots \\ \tilde{Q}(\alpha_1)^{(m-1)} & \dots & \tilde{Q}(\alpha_m)^{(m-1)} \end{pmatrix}. \quad (10)$$

It is well-known that $\tilde{Q}(\alpha_1), \dots, \tilde{Q}(\alpha_m)$ form a basis if and only if the right-hand side of the above equality is a nonsingular matrix and thus if and only if $A(\mathbf{z})$ is nonsingular. In the sequel, one can reduce the problem of the existence of \tilde{Q} satisfying $\text{gcd}(\tilde{Q}, P_{\mathbf{x}=\mathbf{c}}) = 1$ to the problem of the existence of $\mathbf{a} \in k(\mathbf{x})^{\ell}$ in $k(\mathbf{x})$ such that $\det(A(\mathbf{a})) \neq 0$.

Suppose now we already have an operator Q with semisplit coefficients such that P is similar to $P_{\mathbf{x}=\mathbf{c}}$ by Q . Write $Q = \sum_{i=0}^{m-1} b_i \partial_t^i$ where $b_i \in K$ is semisplit. Write further $b_i = \sum_{j=1}^s h_{i,j} \beta_j$ where $h_{i,j} \in k(\mathbf{x})$ and $\beta_j \in k(t) \setminus \{0\}$. Let $L_0 = P_{\mathbf{x}=\mathbf{c}}$ and let \tilde{L}_i be the transformation of L_{i-1} by ∂_t for $i = 1, \dots, m-1$. Then L_i annihilates $\alpha_j^{(i)}$ for all $j = 1, \dots, m$ and $L_i \frac{1}{\beta_l}$ annihilates $\beta_l \alpha_j^{(i)}$ for all $l = 1, \dots, s$ and $j = 1, \dots, m$. Set

$$L = \text{lclm} \left(\left\{ L_i \frac{1}{\beta_l} \mid i = 0, \dots, m-1, l = 1, \dots, s \right\} \right).$$

Then L annihilates all $\tilde{Q}(\alpha_i)$ and thus has P as a right-hand factor. We summarize the previous discussion as the following algorithm.

Algorithm 31 *Input:* $P \in K\langle \partial_t \rangle$ that is monic and completely reducible.

Output: a nonzero $L \in k(t)\langle \partial_t \rangle$ which is divided by P on the right if it exists, otherwise 0.

1. Write

$$P = \partial_t^m + \sum_{i=0}^{m-1} \frac{a_i}{r} \partial_t^i$$

where $a_i \in k(t)[\mathbf{x}]$, $r \in k[\mathbf{x}, t]$.

2. Pick $\mathbf{c} \in k^n$ such that $r(\mathbf{c}) \neq 0$. By the algorithm in [33], compute a basis of the solution space V of the equation (9).

3. Compute a basis of the subspace of V consisting of operators with semisplit coefficients, say Q_1, \dots, Q_ℓ .

4. Set $\tilde{Q} = \sum_{i=1}^\ell z_i Q_i$ and using \tilde{Q} , compute the matrix $A(\mathbf{z})$ as in (10).

5. If $\det(A(\mathbf{z})) = 0$ then return 0 and the algorithm terminates. Otherwise compute $\mathbf{a} = (a_1, \dots, a_\ell) \in k^\ell$ such that $\det(A(\mathbf{a})) \neq 0$.

6. Set b_i to be the coefficient of ∂_t^i in $\sum_{j=1}^\ell a_j Q_j$ and write $b_i = \sum_{j=1}^s h_{i,j} \beta_j$ where $h_{i,j} \in k(\mathbf{x})$ and $\beta_j \in k(t)$. Let $L_0 = P_{\mathbf{x}=\mathbf{c}}$ and for each $i = 1, \dots, m-1$ compute L_i , the transformation of L_{i-1} by ∂_t .

7. Return $\text{lclm} \left(\left\{ L_i \frac{1}{\beta_j} \mid i = 0, \dots, m-1, j = 1, \dots, s \right\} \right)$.

4.2 The general case

Assume that P is (\mathbf{x}, t) -separable and $P = Q_1 Q_2$ where $Q_1, Q_2 \in K\langle\partial_t\rangle$. It is clear that Q_2 is also (\mathbf{x}, t) -separable. One may wonder whether Q_1 is also (\mathbf{x}, t) -separable. The following example shows that Q_1 may not be (\mathbf{x}, t) -separable.

Example 32 *Let $K = k(x, t)$ and let $P = \partial_t^2$. Then P is (\mathbf{x}, t) -separable and*

$$\partial_t^2 = \left(\partial_t + \frac{x}{xt+1} \right) \left(\partial_t - \frac{x}{xt+1} \right).$$

The operator $\partial_t + x/(xt+1)$ is not (\mathbf{x}, t) -separable, because $1/(xt+1)$ is one of its solutions and it is not semisplit.

While, the lemma below shows that if Q_2 is semisplit then Q_1 is also (\mathbf{x}, t) -separable.

Lemma 33 (1) *Assume that $Q_1, Q_2 \in K\langle\partial_t\rangle \setminus \{0\}$, and Q_2 is semisplit. Then $Q_1 Q_2$ is (\mathbf{x}, t) -separable if and only if both Q_1 and Q_2 are (\mathbf{x}, t) -separable.*

(2) *Assume that $P \in K\langle\partial_t\rangle \setminus \{0\}$ and L is a nonzero monic operator in $k(t)\langle\partial_t\rangle$. Then P is (\mathbf{x}, t) -separable if and only if the transformation of P by L is also.*

Proof. Note that the solution space of $\text{lclm}(P_1, P_2)(y) = 0$ is spanned by those of $P_1(y) = 0$ and $P_2(y) = 0$. Hence $\text{lclm}(P_1, P_2)$ is (\mathbf{x}, t) -separable if and only if so are both P_1 and P_2 .

(1) For the “only if” part, one only need to prove that Q_1 is (\mathbf{x}, t) -separable. Assume that g is a solution of $Q_1(y) = 0$ in \mathcal{U} . Let f be a solution of $Q_2(y) = g$ in \mathcal{U} . Such f exists because \mathcal{U} is the universal differential extension of K . Then f is a solution of $Q_1 Q_2(y) = 0$ in \mathcal{U} . By Corollary 20, f is semisplit. Since Q_2 is semisplit, one sees that $g = Q_2(f)$ is semisplit. By Corollary 20 again, Q_1 is (\mathbf{x}, t) -separable.

Now assume that both Q_1 and Q_2 are (\mathbf{x}, t) -separable. Let $\tilde{Q} \in K\langle\partial_t\rangle$ be such that $\tilde{Q} Q_2 = L$ where $L \in k(t)\langle\partial_t\rangle$ is monic. By Corollary 23 and the “only if” part, \tilde{Q} is semisplit and (\mathbf{x}, t) -separable. Thus $\text{lclm}(Q_1, \tilde{Q})$ is (\mathbf{x}, t) -separable. Assume that $\text{lclm}(Q_1, \tilde{Q}) = N\tilde{Q}$ with $N \in K\langle\partial_t\rangle$. Since \tilde{Q} is semisplit, by the “only if” part again, N is (\mathbf{x}, t) -separable. Let $M \in K\langle\partial_t\rangle$ be such that MN is a nonzero operator in $k(t)\langle\partial_t\rangle$. We have that

$$M\text{lclm}(Q_1, \tilde{Q})Q_2 = MN\tilde{Q}Q_2 = MNL \in k(t)\langle\partial_t\rangle.$$

On the other hand, $M\text{lclm}(Q_1, \tilde{Q})Q_2 = M\tilde{M}Q_1Q_2$ for some $\tilde{M} \in K\langle\partial_t\rangle$. Hence $P = Q_1Q_2$ is (\mathbf{x}, t) -separable.

(2) Since L is (\mathbf{x}, t) -separable, we have that P is (\mathbf{x}, t) -separable if and only if $\text{lclm}(P, L)$ is also. Let \tilde{P} be the transformation of P by L . Then $\tilde{P}L = \text{lclm}(P, L)$. As L is semisplit, the assertion then follows from (1). ■

Assume that P is a nonzero operator in $K\langle\partial_t\rangle$. Let P_0 be an irreducible right-hand factor of P . By Algorithm 31, we can decide whether P_0 is (\mathbf{x}, t) -separable or not. Now assume that P_0 is (\mathbf{x}, t) -separable. Then we can compute a nonzero monic operator $L_0 \in k(t)\langle\partial_t\rangle$ having P_0 as a right-hand factor. Let P_1 be the transformation of P by L_0 . Lemma 33 implies that P is (\mathbf{x}, t) -separable if and only if P_1 is also. Note that

$$\begin{aligned} \text{ord}(P_1) &= \text{ord}(\text{lclm}(P, L_0)) - \text{ord}(L_0) \\ &\leq \text{ord}(P) + \text{ord}(L_0) - \text{ord}(P_0) - \text{ord}(L_0) = \text{ord}(P) - \text{ord}(P_0). \end{aligned}$$

In other words, $\text{ord}(P_1) < \text{ord}(P)$. Replacing P by P_1 and repeating the above process yield an algorithm to decide whether P is (\mathbf{x}, t) -separable.

Algorithm 34 *Input: a nonzero monic $P \in K\langle\partial_t\rangle$.*

Output: a nonzero $L \in k(t)\langle\partial_t\rangle$ which is divided by P on the right if it exists, otherwise 0.

1. *If $P = 1$ then return 1 and the algorithm terminates.*
2. *Compute an irreducible right-hand factor P_0 of P by algorithms developed in [4, 32, 34].*
3. *Apply Algorithm 31 to P_0 and let L_0 be the output.*
4. *If $L_0 = 0$ then return 0 and the algorithm terminates. Otherwise compute the transformation of P by L_0 , denoted by P_1 .*
5. *Apply Algorithm 34 to P_1 and let L_1 be the output.*
6. *Return L_1L_0 .*

The termination of the algorithm is obvious. Assume that $L_1 \neq 0$. Then $L_1 = Q_1P_1$ for some $Q_1 \in K\langle\partial_t\rangle$. We have that $P_1L_0 = \text{lclm}(P, L_0)$. Therefore

$$L_1L_0 = Q_1P_1L_0 = Q_1\text{lclm}(P, L_0) = Q_1Q_0P$$

for some $Q_0 \in K\langle\partial_t\rangle$. This proves the correctness of the algorithm.

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